

## HIGH LEVEL COMPARISON - A METHOD OF DETERMINING THE CLASS OF THE DECISION

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## ABSTRACT

This article covers complex modules

$$f(x) \equiv 0 \left( mod \ p_1^{\alpha_1} \ p_2^{\alpha_2} \dots p_k^{\alpha_k} \right)$$

comparisons are output to the comparison system, and the solutions of each comparison of the system are determined using the product.

KEYWORDS. Comparison, system of comparisons, product, remainder.

We have a complex module,  $f(x) \equiv 0 \pmod{p_1^{\alpha_1} p_2^{\alpha_2} \dots p_k^{\alpha_k}}$  (1)give a comparison. Here are  $p_1, p_2, \dots, p_k$  different prime numbers,  $(P_i, P_j) = 1$   $i \neq j$ ,  $i = \overline{1, k}$ ,  $j = \overline{1, k}$ .

(1) Comparisons should be required to define a class of solutions. Usually given

(1) is equivalent to the following system of comparison:

$$\begin{cases} f(x) \equiv 0 \pmod{P_1^{\alpha_1}} \\ f(x) \equiv 0 \pmod{P_2^{\alpha_2}} \\ \dots \\ f(x) \equiv 0 \pmod{P_k^{\alpha_k}} \end{cases}$$
(2)

(2) the relationship is reasonable that is, comparison of property come turns out to be .[1], [2]

Seconds on the other hand, in general, when a high degree of comparison dice universal formula, solving is not. Therefore, the possibility of boric General without  $f(x) \equiv 0 (modP^{\alpha})$  (3) solving the training class let's find out. This method is originally  $f(x) \equiv 0 (modP)$  (4).

Here  $_f(x) = a_0 x^n + a_1 x^{n-1} + \dots + a_n$  (5)

If n > p in (5), this comparison level (p-1) can be reduced.  $f(x) = (x^p - x)Q(x) + R(x)$  (6)



Given the spot  $R(x) \equiv 0 \pmod{P^{\alpha}}$  (5)'in (6), the comparison in (5) can be written as follows: $(x^{p} - x) \neq p$ 

(5) neither initially module P on the classroom solutions let's learn .

Assumption Let (5) — has a solution and should \_ get:

 $x \equiv x_1(modP) \quad \rightarrow x = x_1 + Pt_1 \quad (7)$ 

(7) c (5) ha put her \_ decision 
$$P^2$$
 module on let's find our

 $f(x_1) + Pf'(x_1)t \equiv 0 (modP^2) \qquad f(x_1) \vdots PConsidering$  $\frac{f(x_1)}{P} + f'(x_1)t_1 \equiv 0 (modP) \qquad (8)$ 

(8) to the decision has \_ If  $(f'(x_1), P) = 1$  so.  $x = x_1 + Pt_1$  with the value of R (x) Taylor in a row distribute  $P^2$  ha caralla hadlar leave will be sent.

The last condition from the solution (8) to the solution has as well as this solution class as follows to obtain :

$$t_1 \equiv t'_1(modP) \rightarrow x = x_1 + P(t'_1 + P^2t_2) = x_2 + P^2t_2$$
$$x_2 = x_1 + Pt'_1 \qquad (9)$$

(9) neither (5) ha put the comparison  $P^3$  module on the solutions class we learn and the yield that was ifodada  $P^3$  can be written as follows:  $f(x_2) + f'(x_2)P^2t_2 \equiv 0 \pmod{P^3}$ 

 $\frac{f(x_2)}{P^2} + f'(x_2)t_2 \equiv 0 \pmod{P}$  it's here too  $(f'(x_2), P) = 1$  given that the final solution for the comparison above can be written as follows  $t_2 \equiv t'_2 + Pt_3$  (10) (10) neither (9) ha let's put  $x = x_2 + P^2(t'_2 + Pt_3) = x_2 + P^2t'_2 + P^3t_3$  Repeating this

process  $(\alpha - 2)$ , the class of general solutions can be written as follows: $x = x_{\alpha} + P^{\alpha}t_{\alpha}$ 

The above listed method is the following condition reasonable that again while carra emphasizing let's go  $(f'(x_k, P^k) = 1, k = \overline{1, \alpha})$ 

The article  $f(x) \equiv 0 \pmod{p_1^{\alpha_1} p_2^{\alpha_2} \dots p_k^{\alpha_k}}$  apparently high level comparison to solve the training class comparison to the citation system  $f(x) \equiv 0 \pmod{p_i^{\alpha_i}}$   $i = \overline{1, k}$  then discusses the considerations for solving the decision class of each system of comparisons using the product, and the corresponding

 $f(x) \equiv 7x^3 + 19x + 25 \equiv 0 \pmod{27}$  comparison decision class  $x \equiv 13 \pmod{27}; \quad x = 13 + 27t_3$  found. For example  $f(x) \equiv 0 \pmod{3^3}$  (11)  $f(x) = 7x^3 + 19x + 25$ 

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 $\begin{aligned} f(x) &\equiv 7x^3 + 19x + 25 \equiv x^3 + x + 1 \equiv 0 \pmod{3}, (x^3 - x) + 2x + 1 \equiv 0 \pmod{3} \\ (x^3 - x) &\vdots 3 \to 2x \equiv -1 \pmod{3}, \ 2x \equiv 2 \pmod{3}, \ x \equiv 1 \pmod{3}, \ x = 1 + 3t_1 \end{aligned} (12) \\ (12) \text{ to } (11) \text{ release } \_ 3^2\text{Ha divisible limits leave send} \\ f'(x_1) &= (28x^3 + 19)x \equiv 47, \ f(1) = 51, f(1) + f'(1)3t_1 \equiv 0 \pmod{3^2}, \\ 51 + 3 \cdot 47t_1 \equiv 0 \pmod{3^2}, 17 + 47t_1 \equiv 0 \pmod{3}, \ -2 + 2t_1 \equiv 0 \pmod{3}, \\ t_1 \equiv 1 \pmod{3}, t_1 = 1 + 3t_2 \end{aligned} (13) \\ (13) \text{ to } (12) \text{ let's put} x = 1 + 3(1 + 3t_2) = 4 + 3^2t_2 \end{aligned} (14) \qquad x_2 = 4 \\ f(4) &= 7 \cdot 4^3 + 19 \cdot 4 + 25 = 7 \cdot 64 + 76 + 25 = 549 \\ f'(x) &= 21x^2 + 19, f'(4) = 21 \cdot 4 + 19 = 84 + 19 = 103, 549 + 103 \cdot 3^2t_2 \equiv 0 \pmod{27} \\ &= 61 + 103t_2 \equiv 0 \pmod{3}, \ 1 + t_2 \equiv 0 \pmod{3}, t_2 \equiv 2 + 3t_3 \end{aligned} (15) \end{aligned}$ 

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